



Computer Architecture II - WS 05/06

(due: Monday, 31.10.2005)

Excercise 1: (Number Representation)

(10 points)

In the lecture we defined the binary representation of numbers. We will now extend that definition.

$$B, h \in \mathbb{N}; B \geq 2; 1 \leq h \leq (n-1)$$

Def:

$$\langle a[n-1:0] \rangle_B = \sum_{i=0}^{n-1} a_i \cdot B^i$$

To be proven:

$$\langle a[n-1:0] \rangle_B = \langle a[n-1:h] \rangle_B \cdot B^h + \langle a[h-1:0] \rangle_B$$

Excercise 2: (Recursive Estimations Warmup)

(5 points)

In the lecture we treated the construction of *conditional sum adders*. The following information was given:

$$D(1) = D(FA)$$

$$D(n) = D(Mux) + D\left(\frac{n}{2}\right)$$

Assume that n is a power of two ($n = 2^k$). To be proven:

$$D(n) = \log_2 n \cdot D(Mux) + D(FA)$$

Excercise 3: (Recursive Estimations Continued)

(10 points)

In the lecture we treated the construction of a *compound adder*. We can estimate the cost of the construction using:

$$C(1) \leq 2 \cdot C(FA)$$

$$C(n) \leq C\left(\frac{n}{2}\right) + (n+2) \cdot C(Mux)$$

Assume that n is a power of two ($n = 2^k$). To be proven:

$$C(n) \in \mathcal{O}(n \log_2 n)$$



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Excercise 4: (*Conditional Sum Adder Correctness*)

(15 points)

Prove the correctness of the *conditional sum adder* given in the lecture. You can assume the correctness of a *full adder*:

$$x, y, z, c \in \{0, 1\}, a, b, s \in \{0, 1\}^n$$

$$FA \text{ implements: } \langle c_{out}, z \rangle = \langle x \rangle + \langle y \rangle + c_{in}$$

Assume that n is a power of two ($n = 2^k$). To be proven:

$$COS_n \text{ implements: } \langle a \rangle + \langle b \rangle + c_{in} = \langle c_n, s[n-1:0] \rangle$$